## Seminar Logic and Foundations of Computing Homework 3

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**Problem 1.** Show that every colimit can be expressed as a filtered colimit of finite (finitely generated) colimits, viz. given a diagram  $D : \mathcal{D} \to \mathcal{C}$  with  $\mathcal{C}$  cocomplete and  $\mathcal{D}$  small, colim D can be constructed as the filtered colimit of the diagram of all colim D', where  $D' : \mathcal{D}' \to \mathcal{D}$  ranges over all domain restrictions of  $\mathcal{D}$  to finitely generated subcategories  $\mathcal{D}'$  of  $\mathcal{D}$ .

Solution. Let  $\mathcal{C}$  be cocomplete, and let  $D: \mathcal{D} \to \mathcal{C}$  be a small diagram. Then we have a diagram  $F : \operatorname{FinSub}(\mathcal{D}) \to \mathcal{C}$  mapping finitely generated subcategories  $\mathcal{D}'$  of  $\mathcal{D}$  to their respective colimits, and mapping an inclusion morphism  $f: \mathcal{D}' \to \mathcal{D}''$  to the factorisation of colim D'' through colim D'. Then we have to show a few things: We have to show that the category  $\operatorname{FinSub}(\mathcal{D})$  is filtered, and we need to show that indeed these colimits are the same.

- 1. The category of finite subcategories is filtered. We can see this by using Lemma 2.19  $(3 \implies 1)$ . Firstly, note that the empty category is finitely generated, and therefore a subcategory of  $\mathcal{D}$ , Meaning that FinSub( $\mathcal{D}$ ) is nonempty. Secondly, every pair of finitely generated subcategories  $\mathcal{D}', \mathcal{D}''$  has a span given by the subcategory generated by the union of their objects and the union of their morphisms. Then clearly,  $\mathcal{D}'$  and  $\mathcal{D}''$  are subcategories, meaning we have inclusion morphisms which are a span. Thirdly, when we have two parallel morphisms  $\mathcal{D}' \xrightarrow{f,g} \mathcal{D}''$ , we have a morphism equalizing them: The identity on  $\mathcal{D}''$ , since inclusions are identical in this category.
- 2. To show that the colimits are the same, we will show that either colimit forms a cone over the others diagram. First, for each finitely generated subdiagram D' of D, clearly we have a cone of the colimit of D over D'. As a result, we find that there are maps  $\operatorname{colim}_{d\in \mathcal{D}'} D' \to \operatorname{colim}_{d\in \mathcal{D}} D$  which factor all maps  $p_d$  (see diagram below) for  $d \in \mathcal{D}'$ . This creates a cone over F, since for any inclusion in FinSub $(\mathcal{D})$ , our choice of morphism must commute by virtue of forming a cone over the full diagram of D. As a result, we find that there is a unique map from colim F to colim D.



Second, we will show that the colimit of F forms a cone over D: for every object  $d \in \mathcal{D}$ , form the finitely generated subcategory of  $\mathcal{D}$  containing just that object. Since clearly the colimit of this diagram is equivalent to itself, we find that  $\rho_d : Dd \to \operatorname{colim} F$  (up to isomorphism). This is indeed a cone, since for any morphism  $f : d \to d'$ , we have inclusion maps from the finitely generated subcategories  $\{d\}$  and  $\{d'\}$  to the finitely generated subcategory corresponding with the graph  $\{d \xrightarrow{f} d'\}$ . These inclusion maps give rise to morphisms from  $\operatorname{colim} \{Dd\}$  and  $\operatorname{colim} \{Dd'\}$  to  $\operatorname{colim} \{Dd \xrightarrow{Df} Dd'\}$  which commute with the injections  $\rho_{\{Dd\}}, \rho_{\{Dd'\}}$  and  $\rho_{\{Dd \xrightarrow{Df} Dd'\}}$ , proving that indeed the morphisms out of the singleton diagrams form a cone over D, giving us a map  $\operatorname{colim} D \to \operatorname{colim} F$ .

Finally, since the mappings of cones compose to precisely the identity on both sides, we find that their respective factorings must form an equivalence, completing our proof.  $\Box$